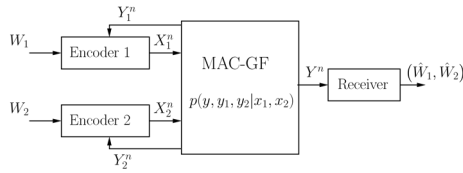


Motivation

- Inherent openness of wireless medium brings two concepts:
 - Cooperation:** Overheard information allows users to cooperate and increases rates
 - Secrecy:** But it also decreases secrecy
- How do cooperation and secrecy interact?
- Our focus:
 - A two user multiple access channel with generalized feedback
 - Each user considers the other as a wire-tapper
 - Yet, they are still willing to cooperate through feedback
 - How does **cooperation** affect **secrecy**?

* Feedback increases achievable rates, but what about secrecy?



MAC with Generalized Feedback

- Presence of feedback: **Cooperation** and/or **Secrecy**
- Cooperation:
 - Similar to the relay channel, basic schemes are decode-and-forward (DAF) and compress-and-forward (CAF) [Cover-El Gamal 1979]
 - DAF: Cover-Leung 1981, Willems *et al.* 1983, Carleial 1982
 - CAF: Khojastepour 2004, Ong-Motani 2007
- Secrecy:
 - Previous works did not utilize the feedback, i.e., no cooperation
 - The only effect of feedback: Leakage of confidential information
 - * Liang-Poor 2008, Liu-Maric-Yates-Spaorojevic 2006

Generalized MAC with Secrecy

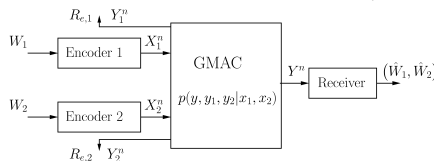


Figure 1: The channel model in Liang-Poor 2008, Liu *et al.* 2006.

An Achievable Scheme

Theorem 1 Rate tuples $(R_1, R_2, R_{e,1}, R_{e,2})$ satisfying

$$\begin{aligned} R_1 &\leq R'_1 && \leq I(X_1; Y, \hat{Y}_1 | U, X_2) \\ R_2 &\leq R'_2 && \leq I(X_2; Y, \hat{Y}_1 | U, X_1) \\ R_1 + R_2 &\leq R'_1 + R'_2 && \leq I(X_1, X_2; Y, \hat{Y}_1 | U) \\ R_{e,1} &\leq \min \{ R'_1 - \tilde{R}_{e,1}, R_1 \} \\ R_{e,2} &\leq \min \{ R'_2 - I(X_2; Y_1 | U, X_1), R_2 \} \end{aligned}$$

where $\tilde{R}_{e,1}$ is given by

$$\tilde{R}_{e,1} = \begin{cases} I(X_1; Y_2, \hat{Y}_1 | U, X_2) & \text{if } s_1 \\ I(X_1; Y_2 | U, X_2) & \text{otherwise} \end{cases}$$

are achievable for any distribution of the form

$$p(u)p(x_1|u)p(\hat{y}_1|u, x_1, y_1)p(x_2)p(y, y_1, y_2|x_1, x_2)$$

subject to the constraint

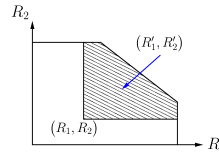
$$I(\hat{Y}_1; Y_1 | U, X_1) \leq I(U, \hat{Y}_1; Y)$$

where

$$s_1 = \{I(U; Y) \leq I(U; Y_2|X_2), \quad I(\hat{Y}_1; Y|U) \leq I(\hat{Y}_1; Y_2|U, X_2)\}$$

Comments on the Achievable Scheme

- It uses CAF for cooperation
- Stochastic encoding:
 - Embed your messages (R_1, R_2) to a larger codebook (R'_1, R'_2)
 - Additional redundancy is to confuse the other user



- User 1's observation can include self-interference
 - We eliminate user 1's signal, X_1 , from Y_1 while compressing it to \hat{Y}_1
 - An alternative: Keep self-interference for a possible rate increase
- Channel prefixing can be employed, replace X_1 (resp. X_2) with V_1 (resp. V_2) and the PMF with

$$p(u)p(v_1|u)p(x_1|v_1)p(\hat{y}_1|u, v_1, y_1)p(v_2)p(x_2|v_2)p(y, y_1, y_2|x_1, x_2)$$
- User 2 may want to decode the compressed version of user 1's feedback, \hat{Y}_1
 - Depending on whether it can or not, different leakage expressions
- We can also extend this scheme to the case with two-sided cooperation [Ekrem-Ulukus CISS 2008]

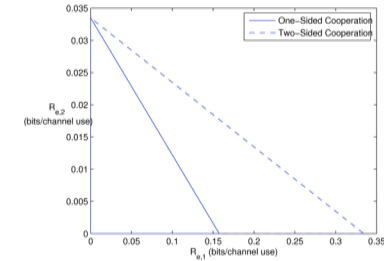
Gaussian Channels

- A Gaussian MAC-GF is

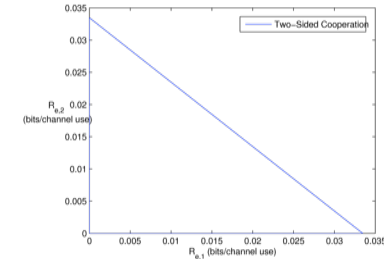
$$\begin{aligned} Y_i &= X_1 + X_2 + Z_i, \quad i = 1, 2 \\ Y &= X_1 + X_2 + Z \end{aligned}$$

where $Z \sim \mathcal{N}(0, N)$, $Z_i \sim \mathcal{N}(0, N_i)$, $i = 1, 2$. $E[X_1^2] \leq P_1$ and $E[X_2^2] \leq P_2$.

- For Gaussian channels without cooperation,
 - If $N_1 < N$, then $R_{e,2} = 0$ and if $N_2 < N$, then $R_{e,1} = 0$
- What is the effect of cooperation on secrecy?
 - Case I:** No secrecy for user 2 if cooperation is not allowed, $N_1 = 0.75 < N = 1 < N_2 = 1.25$



- Case II:** No secrecy for both users if cooperation is not allowed, $N_1 = N_2 = 0.75 < N = 1$



Conclusions

- Cooperation can increase secrecy; even an untrusted party can help
- Since cooperation enlarges achievable rates, this might be expected
- Crucial point: Depends on how cooperation enlarges the achievable region
 - DAF or partial DAF cannot increase secrecy, but CAF can.
 - Via CAF, a user can increase the rate of the other user beyond its own decoding capability.