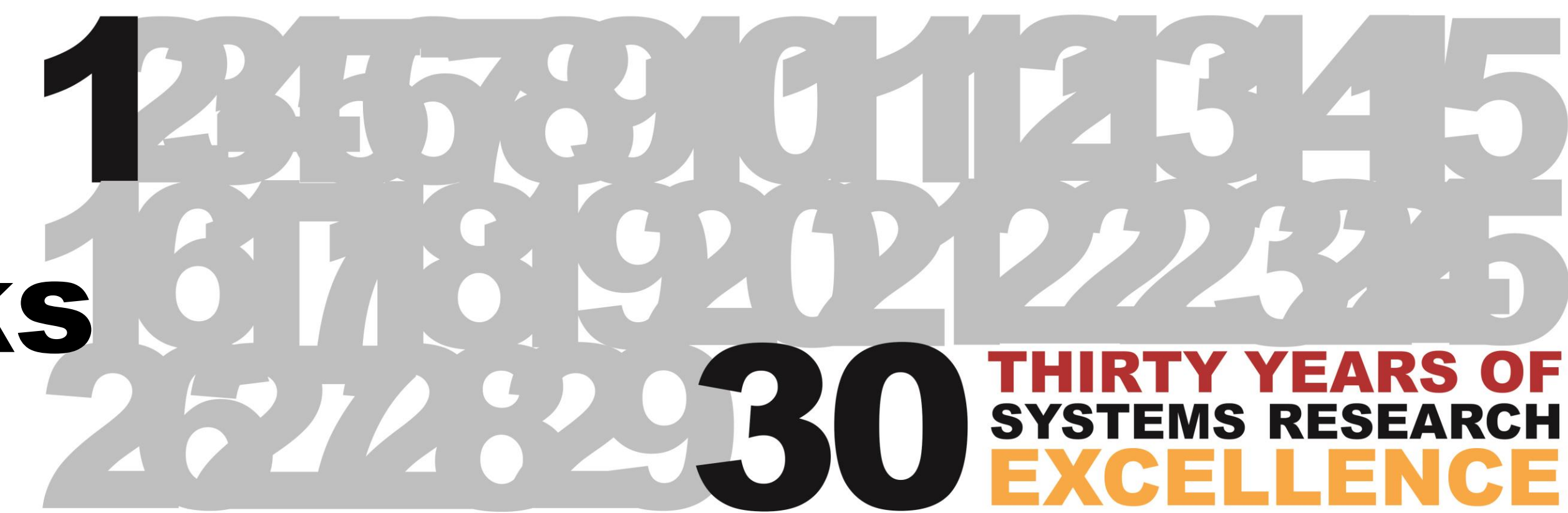


Mathematical Programming Models for Influence Maximization on Social Networks

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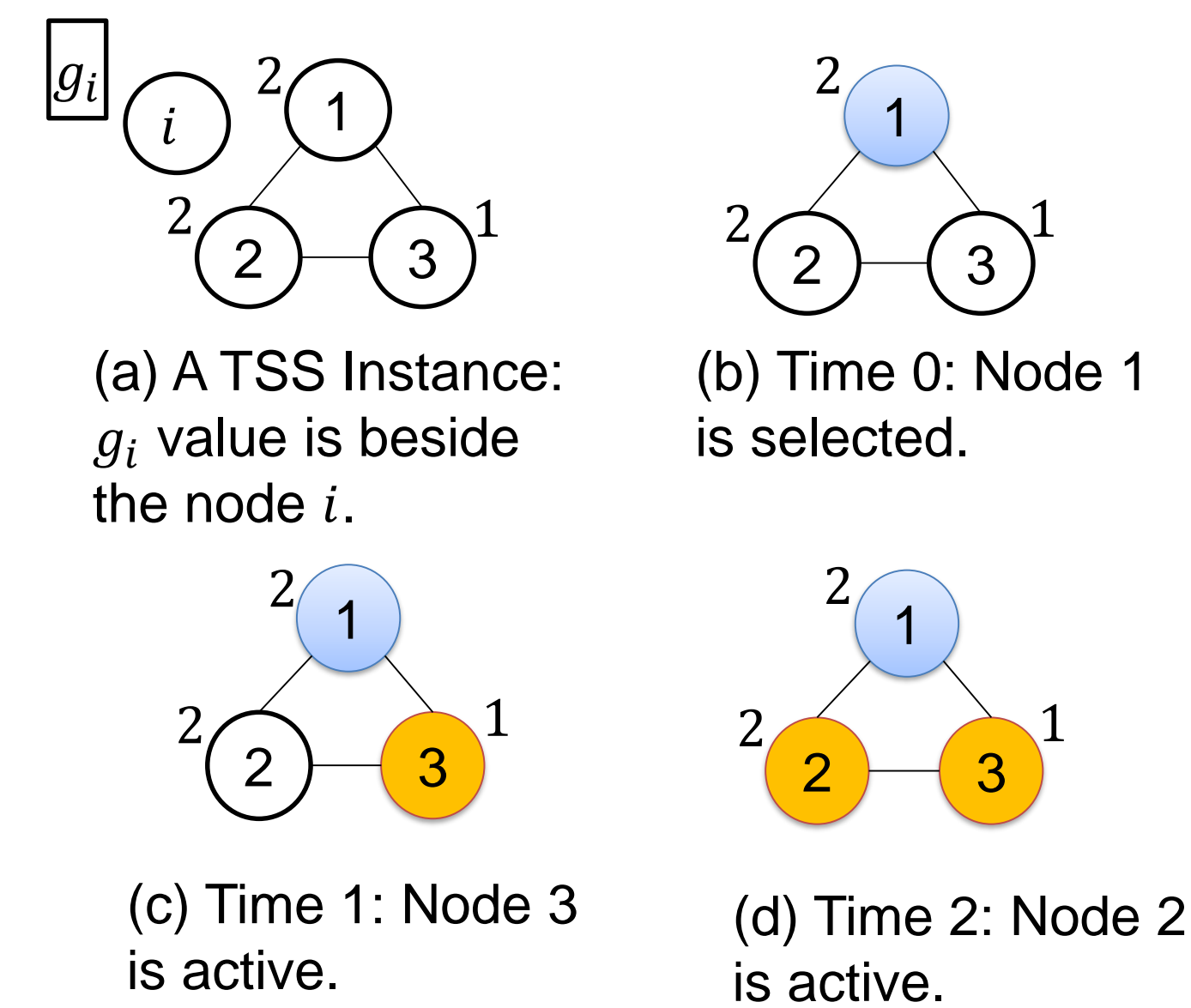
History of the Research



Fundamental problem in social network analysis: Whom do you target to maximize the adoption of a product/innovation? Also has applications in epidemiology.

- Chen [2009] proposed the **Target Set Selection (TSS)** problem:
- Given a connected undirected graph $G = (V; E)$. For each $i \in V$, there is a threshold, g_i , which is between 1 and $\text{degree}(i)$. All nodes are inactive initially.
 - Select a subset of nodes, the target set, and they become active immediately.
 - After that, in each step, an inactive node i becomes active if at least g_i of its neighbors are active in the previous step.
 - Goal: Find the minimum target set while ensuring that all nodes are active at the end of this diffusion process.**

All previous work has focused on heuristics and approximation algorithms.



We want to apply mathematical programming techniques to develop **EXACT** approaches.

Future of the Research

First researchers to study exact solution approaches to problem and to solve large real world size instances to optimality! Focus of future research:

- Proportion requirements---fraction α of population is influenced (instead of 100% adoption).
- Latency constraints---time constraints on diffusion process.
- Time varying networks---connections change over time.
- Robust target sets/variants under stochastic influence factors.
- Applications in epidemiology.
- Voting models and opinion formation on social networks.

Current State of the Research

We study three fundamental problems. All three problems are NP-hard (APX-hard).

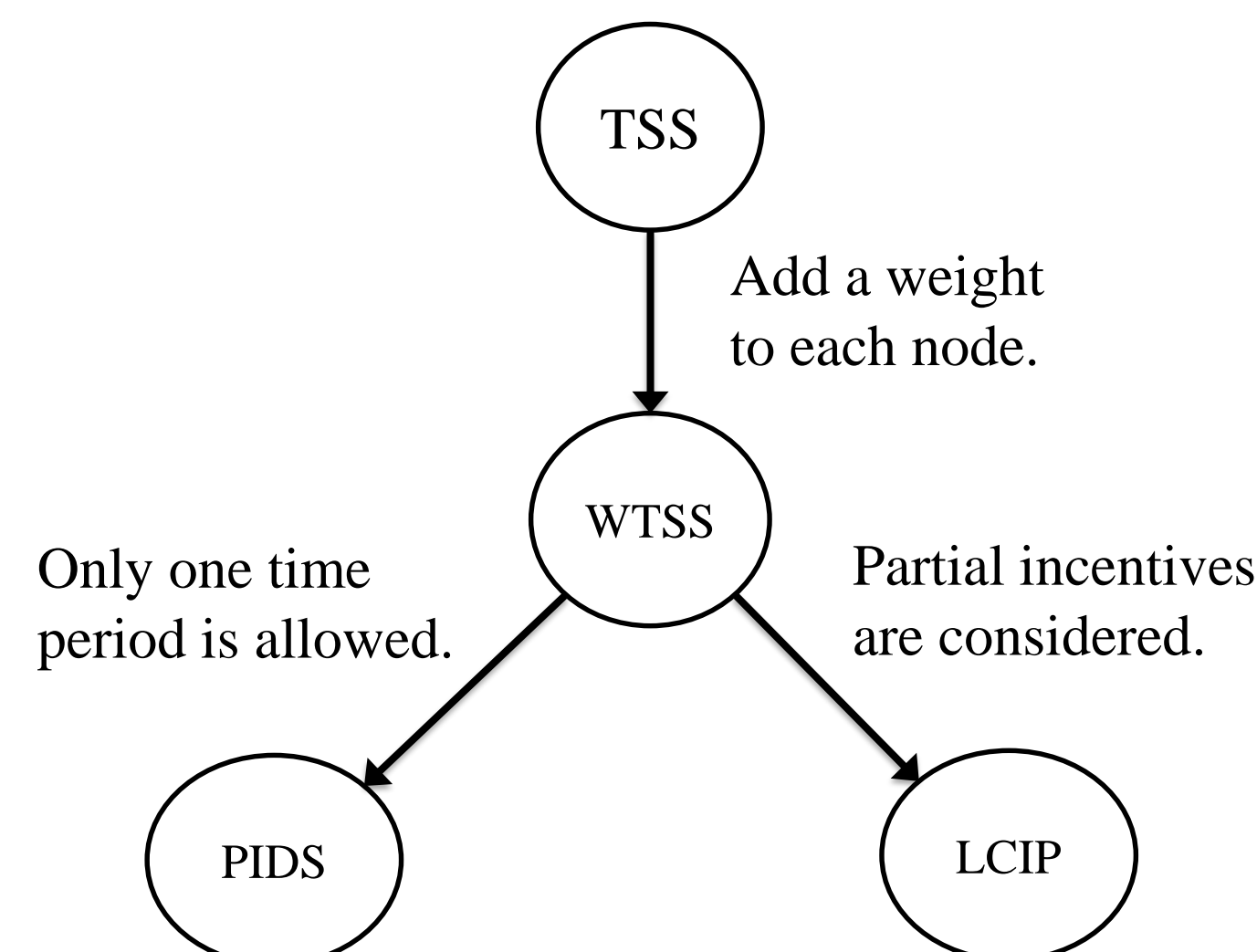


Fig: Relationships among these three problems.

Our Research Scheme



Start with tree graphs:

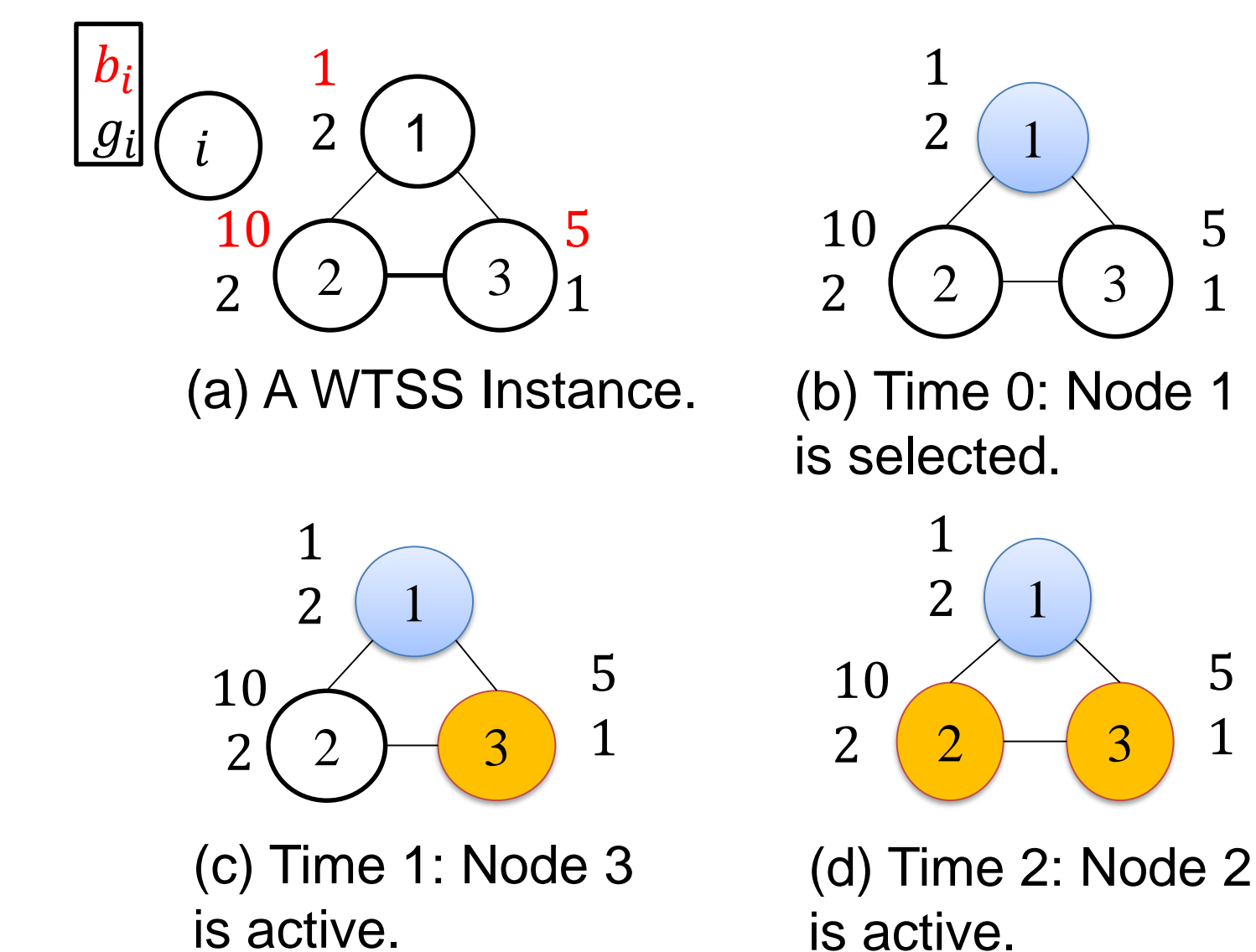
- Can we solve it?
- Can we find the strongest integer programming (IP) formulation for it?
- Can we study its polytope?

Turn to general graphs:

- How to apply the "good" tree formulation here?
- Can we develop Branch-and-Cut approaches?

The Weighted Target Set Selection (WTSS) Problem

A weight (cost) b_i for each node $i \in V$ (different nodes require different levels of effort).

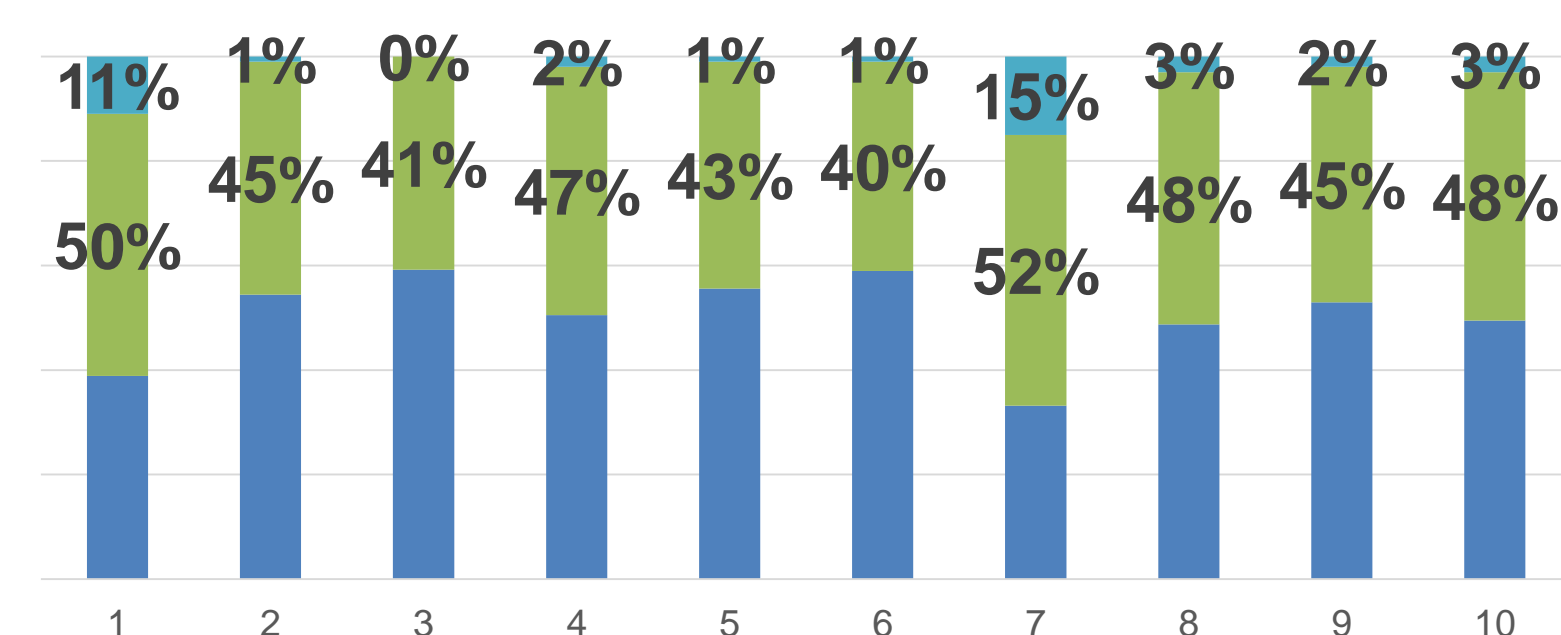


Our results for the WTSS and PIDS problems:

- On trees:**
- A linear time algorithm.
 - A tight and compact extended IP formulation.
 - A complete description of the polytope.
- On general graphs:**
- A strong IP formulation.
 - A specialized Branch-and-Cut approach.

Some computational results for the WTSS:

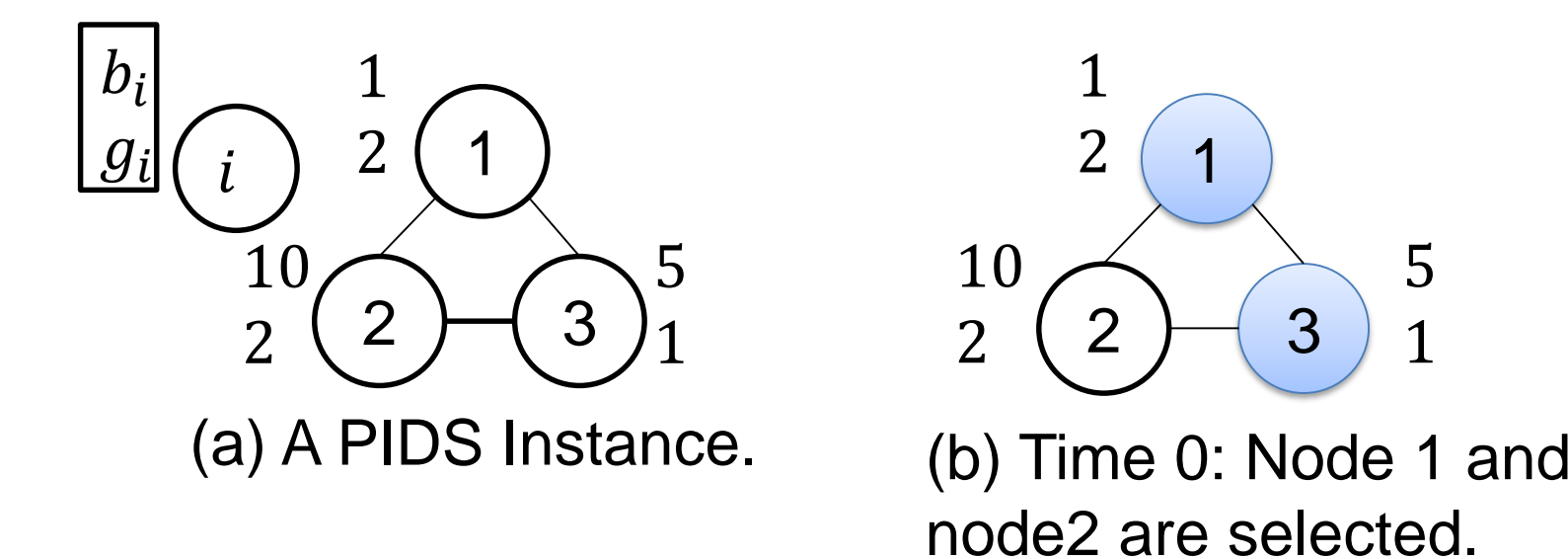
- 1000 nodes, 4000 edges, 10 instances.



Blue: previous formulation. Green: gap we closed. Aqua: the remaining gap.

The Positive Influence Dominating Set (PIDS) Problem

Only **one** time period is allowed for diffusion.

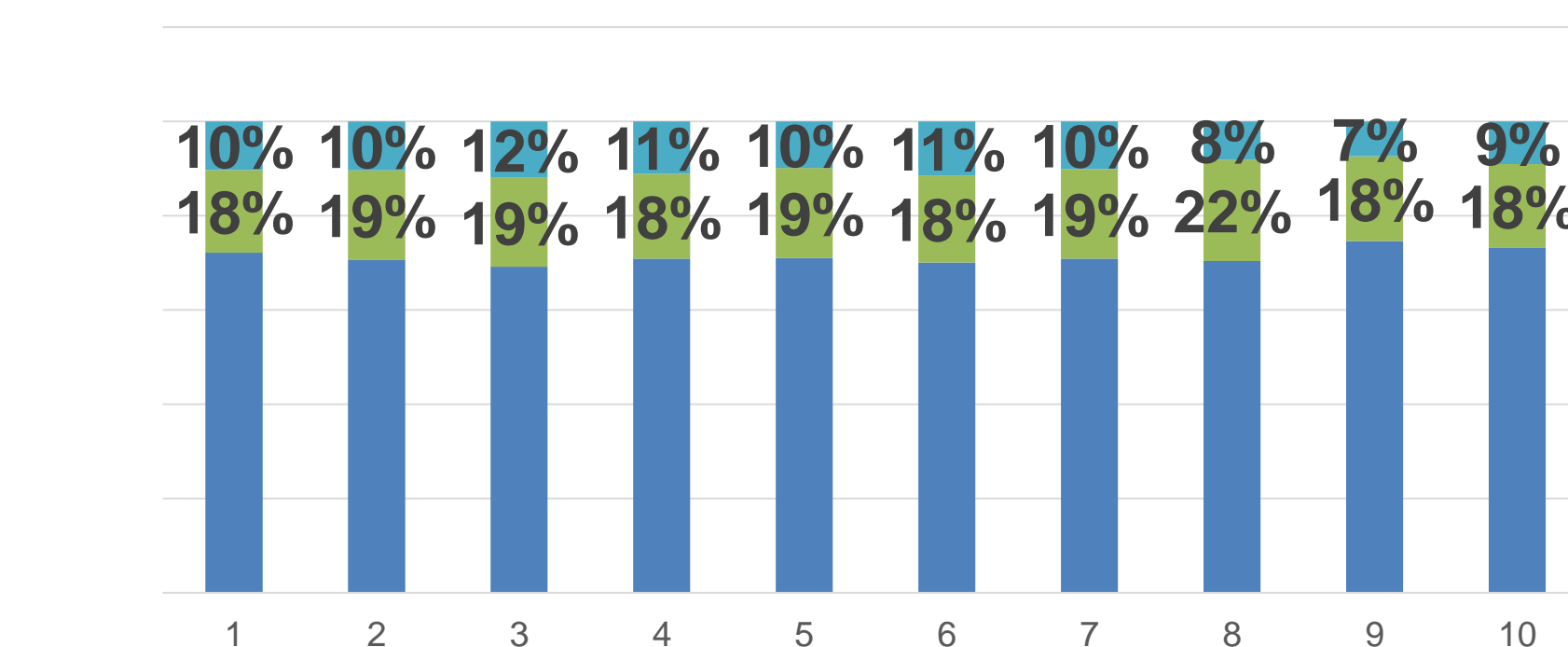


Instances and Environment:

- Generate networks à la Watts and Strogatz [1998] model.
- Randomly generate b_i and g_i for a node i .
- CPLEX 12.6, Python API, Intel i5 3.40GHz, 24 GB ram, Ubuntu.

Some computational results for the PIDS problem:

- 200 nodes, 800 edges, 10 instances.



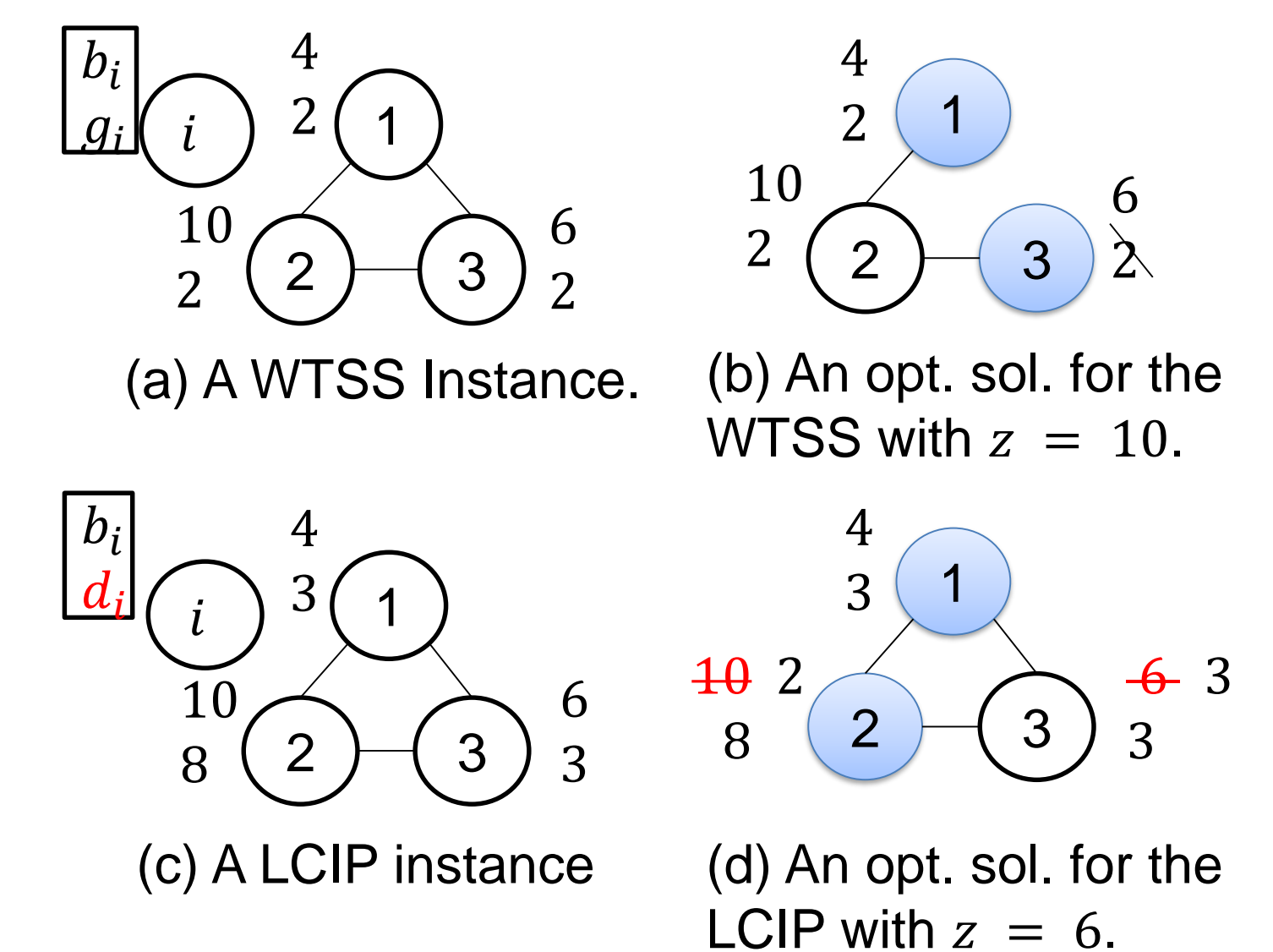
Blue: previous formulation. Green: gap we closed. Aqua: the remaining gap.

Our results for the LCIP:

- A totally unimodular matrix (TUM) formulation on trees.
- A Branch-and-Cut approach for general graphs based on the TUM formulation and the observation that the influence propagation graph must be a directed acyclic graph.

The Least Cost Influence Problem (LCIP)

- Influence factor d_{ij} represents the influence from node j to node i such that its threshold reduced by d_{ij} after node j becomes active.
- In the WTSS, $d_{ij} = d_i$ and $g_i = \lfloor \frac{b_i}{g_i} \rfloor$.
- In the LCIP, **partial incentives are considered**.



Some computational results for the LCIP:

Tricks of the Branch-and-Cut approach:

- BC**: Branch-and-Cut with cycle separation.
- H**: heuristics for initial solutions.
- C**: prioritize branching over separation.
- B**: a specialized branching rule.
- P**: perturbations for symmetry elimination.

10 **Facebook** graph instances with 4039 nodes and 88234 edges run for 1 hour limit.

	Opt. #	Opt. Run. Time (S)	Avg. Gap
BC-H	0	NA	36.7%
BC-H-C	2	1537	5.8%
BC-H-C-B	2	2535	3.8%
BC-H-C-B-P	5	2271	0.4%